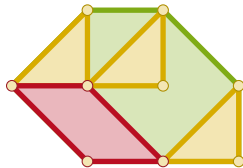


Ehrhart Theory

Polyhedra

▷ $X, Y \subset \mathbb{R}^d$ sets

▶ **Minkowski sum** $X + Y = \{x + y \mid x \in X, y \in Y\}$

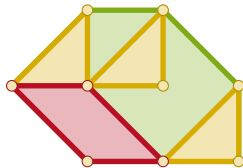


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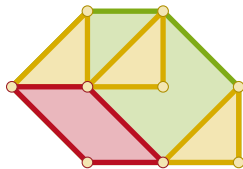
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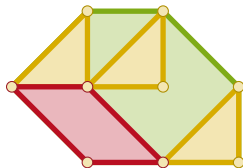
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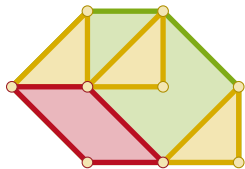
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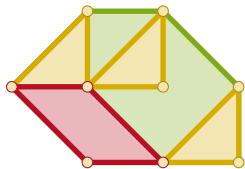
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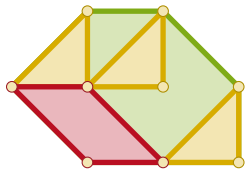
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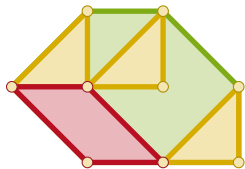
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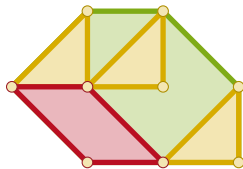
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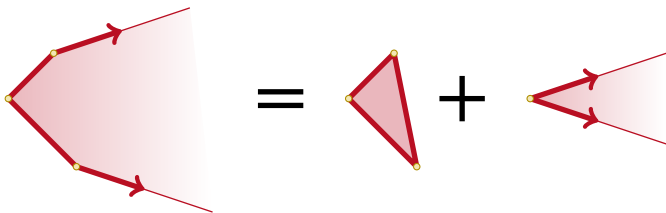
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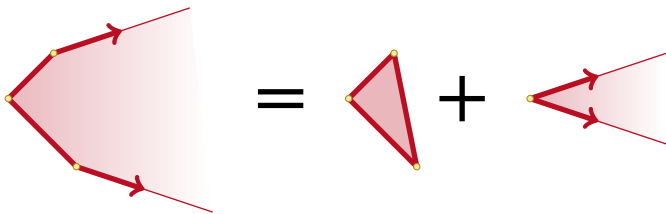
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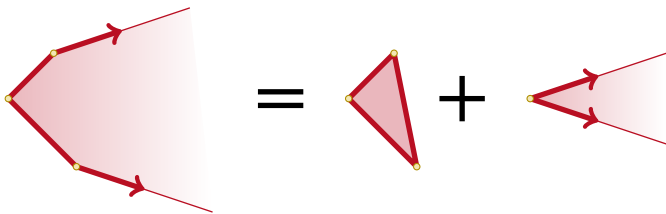
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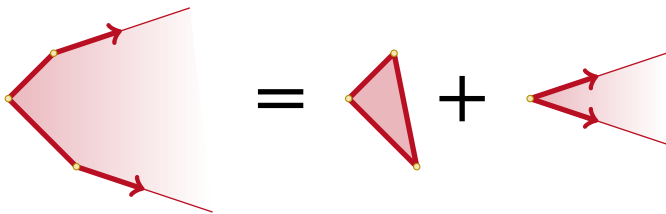
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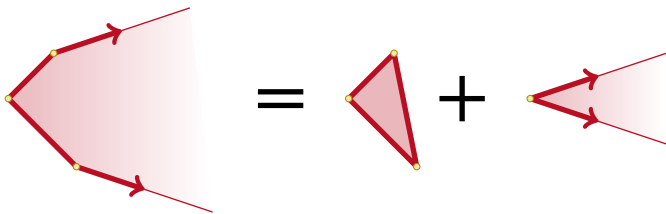
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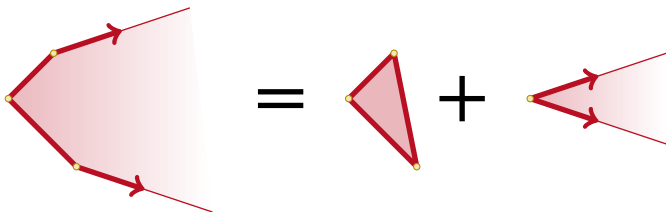
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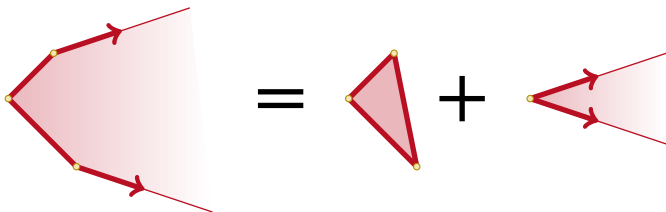
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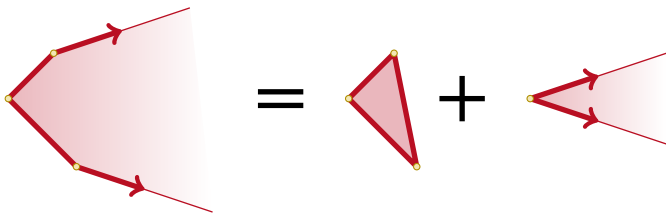
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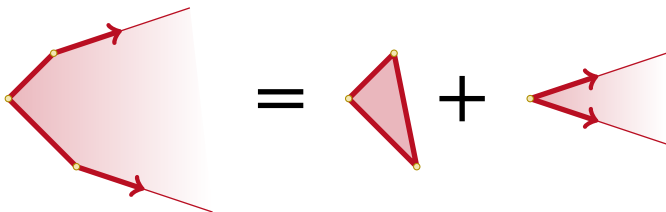
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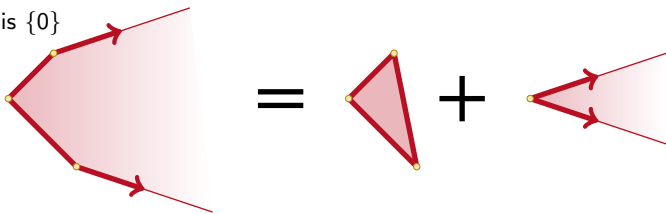
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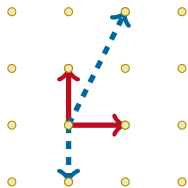
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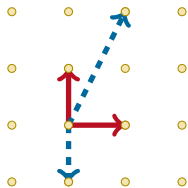
Lattice Polytopes

► $\Lambda \subseteq \mathbb{R}^d$ **lattice** $:\Leftrightarrow$ Λ is discrete, additive, abelian group



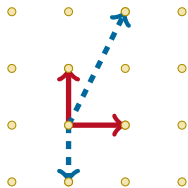
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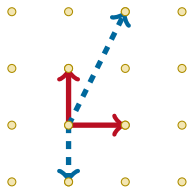
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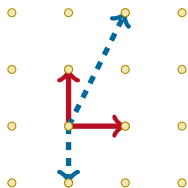
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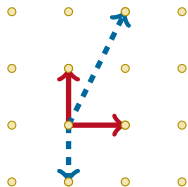
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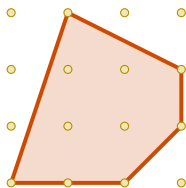
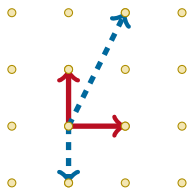
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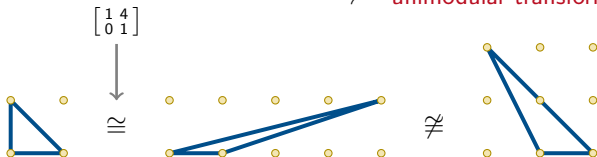
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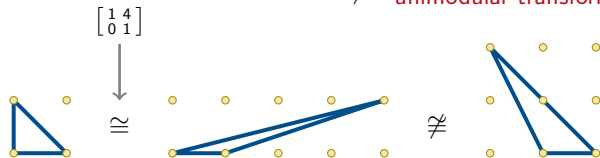
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▶ **Proposition:** unimodular transformations preserve $\#$ (lattice points) and volume (they do *not* preserve lengths and angles)

Ex

Why Count Lattice Points

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- ▶ # ways to pay 2€ with coins of 2€, 1€, 50ct (sufficient supply of coins)?

Why Count Lattice Points

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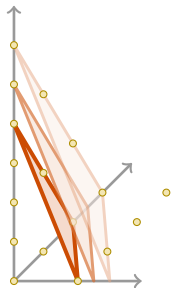
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→ lattice points in the polytope defined by

$$0 \leq x_1, x_2, x_3,$$

$$k = x_1 \cdot 2 + x_2 \cdot 1 + x_3 \cdot \frac{1}{2}$$



Why Count Lattice Points

► contingency tables

	Diploma	PhD	Teacher	
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→ 714,574,663,432 tables.

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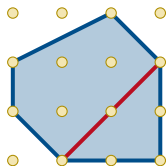
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$$a_1 = 5, \quad i_1 = 3, \quad b_1 = 6$$

$$a_2 = 2, \quad i_2 = 0, \quad b_2 = 6$$

$$e = 1$$

$$a = 5 + 2, \quad i = 3 + 1, \quad b = 6 + 6 - 1 - 1$$

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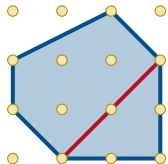
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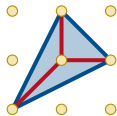
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▷ $b = 3, i \geq 1$: cone over interior point



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▶ further observations: ▷ leading coefficient is volume

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Encoding Lattice Points in Polytopes

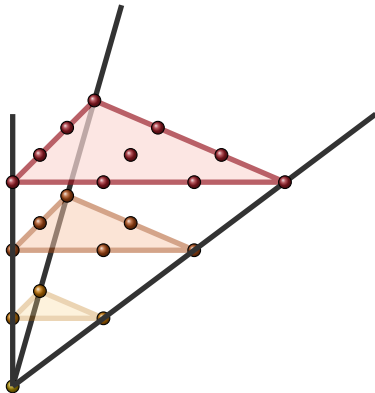
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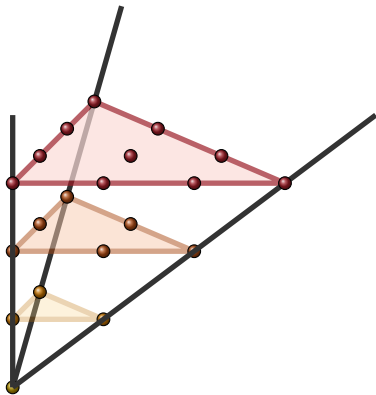
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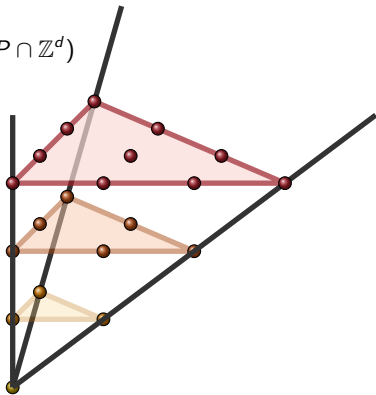
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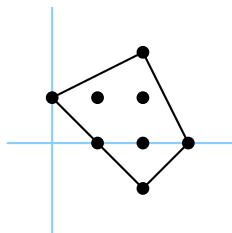
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► example

$$P = \text{conv} \begin{bmatrix} 0 & 2 & 2 & 3 \\ 1 & -1 & 2 & 0 \end{bmatrix}$$



$$\widehat{G}_P := \sum_{a \in P \cap \mathbb{Z}^d} t^a = t_1^2 t_2^2 + t_2 + t_1 t_2 + t_1^2 t_2 + t_1 + t_1^2 + t_1^3 + t_1^2 t_2^{-1} \in \widehat{L}$$

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Proposition: There is a natural homomorphism

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that maps \hat{G} to $\frac{f}{g}$ if $g\hat{G} = f \in L$.

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Encoding Lattice Points in Polytopes

► how do we obtain a rational function?

► **Def:** $\hat{G} \in \hat{L}$ is **summable** $:\iff$ there is $g \in L$ s.th. $g\hat{G} \in L$
 L^{sum} := set of **summable Laurent series**

Example: $\triangleright P = [0, \infty) \longrightarrow \hat{G}(t) = 1 + t + t^2 + \dots = \sum_{a \in \mathbb{Z}_{\geq 0}} t^a$
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► **Def:** P polyhedron, $\hat{G}_P(t)$ summable.

Then $G_P(t) := \Phi(\hat{G}_P)(t)$ is the **integer point generating series**

Generating Functions

▷ $\mathbf{v}_1, \dots, \mathbf{v}_d \in \mathbb{Z}^d$ linearly independent, and $C := \text{cone}(\mathbf{v}_1, \dots, \mathbf{v}_d)$

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→ we have encoded lattice points in a simplicial cone!

Ehrhart Series

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- ▷ Recall the Ehrhart counting function: $\text{ehr}_S(k) := \#(kS \cap \mathbb{Z}^d)$

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Ehrhart Series

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\longrightarrow **non-negative and integral**

Ehrhart's Theorem

► Theorem (Ehrhart's Theorem for Simplices)

S simplex, then

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Ex

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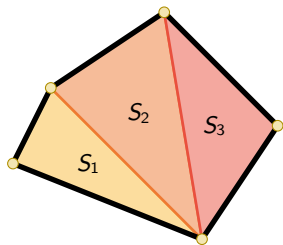
Ehrhart's Theorem

- ▶ how do we extend this to general polytopes?

Ehrhart's Theorem

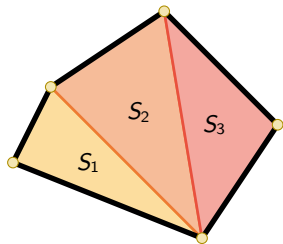
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→ triangulate!



Ehrhart's Theorem

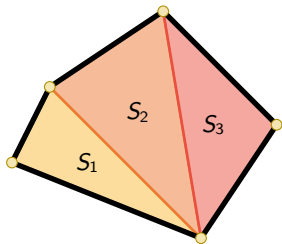
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P lattice polytope

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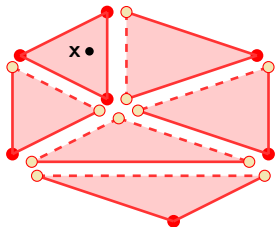
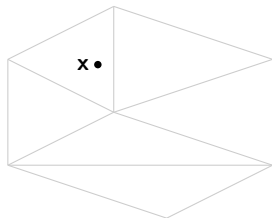
Half-Open Decompositions

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Half-Open Decompositions

► **inclusion-exclusion** is not efficient and difficult to use, loose non-negativity of h_i^*

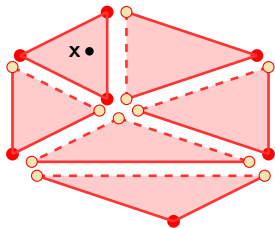
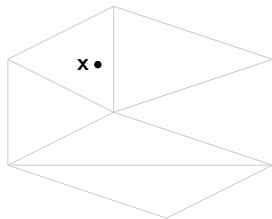
→ do **half-open decompositions**!



Half-Open Decompositions

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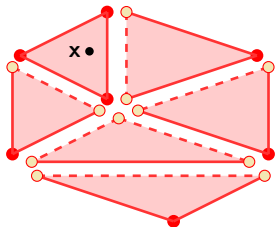
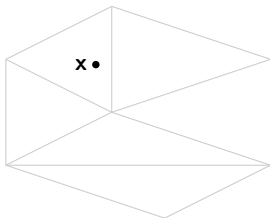


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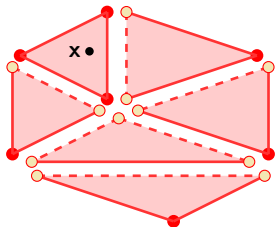
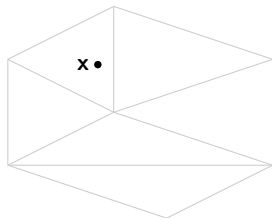


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Half-Open Decompositions

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→ do half-open decompositions!



- ▷ lattice points in intersections only counted once
- ▷ preserves non-negativity
- ▷ need to define this for the cones obtained from such a triangulation

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▷ $C = \text{cone}(V)$ for $V = \{\mathbf{v}_0, \dots, \mathbf{v}_d\}$ simplicial cone in \mathbb{R}^{d+1}

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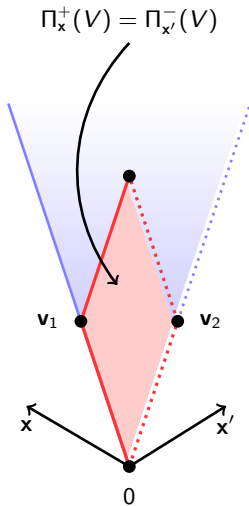
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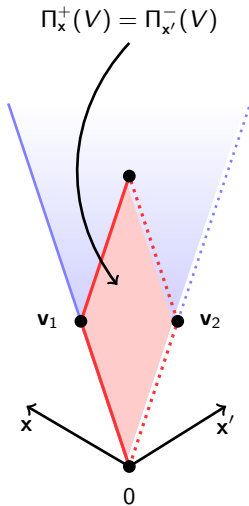
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▷ obtain normal fundamental parallelepiped for $\mathbf{x} \in \text{relint } C$



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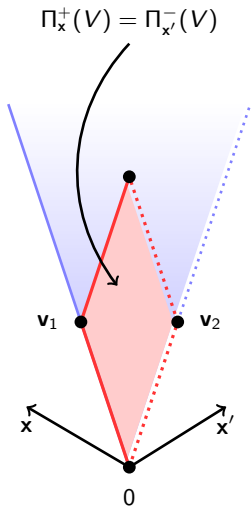
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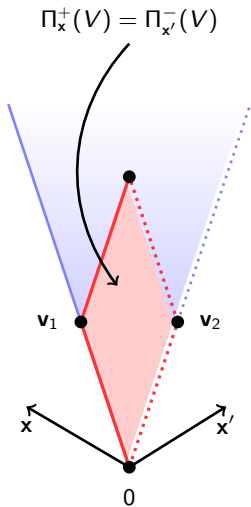
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▶ Prop: $G_{C_{\mathbf{x}}^\pm}(\mathbf{t}) = \frac{G_{\Pi_{\mathbf{x}}^\pm}(\mathbf{t})}{(1 - \mathbf{t}^{\mathbf{v}_0})(1 - \mathbf{t}^{\mathbf{v}_1}) \dots (1 - \mathbf{t}^{\mathbf{v}_d})}$



Ex

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- ▷ D general cone with triangulation \mathcal{T} ,
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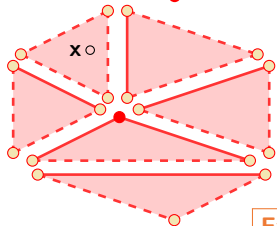
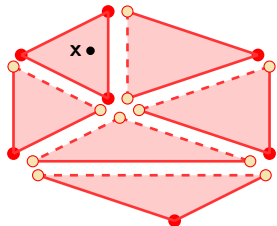
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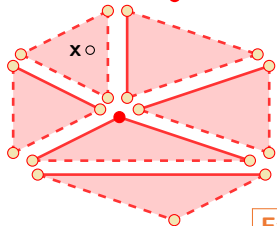
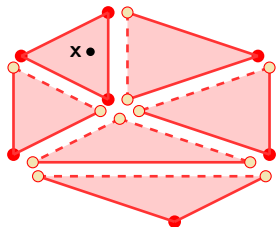
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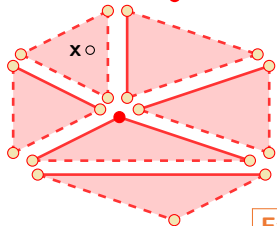
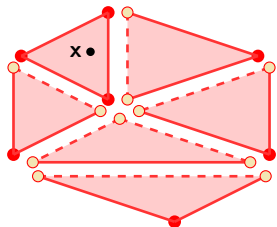


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▷ $G_D(\mathbf{t}) = \sum_{C \in \mathcal{T}[d]} \frac{G_{\Pi_{\mathbf{x}}^+(V(C))}(\mathbf{t})}{\prod_{v \in V(C)} (1 - \mathbf{t}^v)}$ and $G_{\text{relint } D}(\mathbf{t}) = \sum_{C \in \mathcal{T}[d]} \frac{G_{\Pi_{\mathbf{x}}^-(V(C))}(\mathbf{t})}{\prod_{v \in V(C)} (1 - \mathbf{t}^v)}$

Ex

Half-Open Decompositions

- ▷ P lattice polytope with cone $C(P)$
- ▷ \mathcal{T} triangulation of P
 - triangulation **extends** to $C(P)$
- ▷ $\mathbf{x} \in \mathbb{R}^d$ generic with respect to all cones in \mathcal{T}

▶ **Prop:**

$$\text{Ehr}_P(t) = G_{C(P)}(t, \mathbf{1}) = \sum_{S \in \mathcal{T}[d+1]} G_{S_x^+}(t, \mathbf{1}) = \frac{\sum_{S \in \mathcal{T}[d+1]} G_{\Pi(S)_x^+}(t, \mathbf{1})}{(1-t)^{d+1}}$$

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Non-Negativity

► Stanley's Non-negativity Theorem

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Then $h_0^*, \dots, h_d^* \geq 0$ and $h_0^* = 1$

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► There is **no** similar result for the Ehrhart polynomial

► e.g. $S := \text{conv}(0, \mathbf{e}_1, \mathbf{e}_2, \mathbf{e}_1 + \mathbf{e}_2 + 18\mathbf{e}_3)$ has $\text{ehr}_S(t) = 1 - t + t^2 + 3t^3$

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C d -dimensional polyhedral cone with rational generators. Then

$$G_C(\mathbf{t}) = (-1)^d G_{\text{int } C}\left(\frac{1}{\mathbf{t}}\right) \quad \left(\frac{1}{\mathbf{t}} = \left(\frac{1}{t_1}, \dots, \frac{1}{t_d}\right)\right)$$

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► plug this into

$$G_C(\mathbf{t}) = \sum_{S \in \mathcal{T}[d]} G_{S_{\mathbf{x}}^+}(\mathbf{t}) = \sum_{S \in \mathcal{T}[d]} \frac{G_{\Pi_{\mathbf{x}}^+(V)}(\mathbf{t})}{\prod_{\mathbf{v} \in V} (1 - \mathbf{t}^{\mathbf{v}})}$$

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► compare coefficients

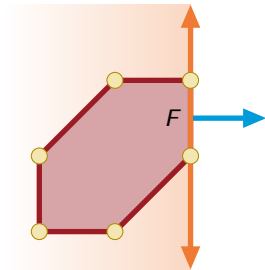
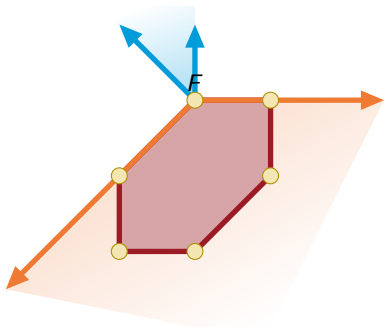
Brion's Theorem

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- ▶ P rational d -dimensional polytope
- ▶ F face of P , the **tangential cone** of F in P is

$$T_F P := \{ \mathbf{x} \mid \exists \mathbf{p} \in F, \varepsilon > 0 : \mathbf{p} + \varepsilon(\mathbf{x} - \mathbf{p}) \in P \}$$

→ up to a shift this cone is dual to the normal cone at this face



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► Theorem (Brianchon-Gram Identity)

P rational d -polytope, then

$$\hat{G}_P(\mathbf{t}) = \sum_{F \leq P} (-1)^{\dim F} \hat{G}_{T_F P}(\mathbf{t})$$

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coefficient on RHS: let $f_i := \#(i\text{-dimensional faces})$

\mathbf{m} is in every target cone, so coefficient is

$$\sum_{F \leq P} (-1)^{\dim F} = \sum_{i=0}^d (-1)^i f_i = 1$$

the last equality follows from Euler's identity

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coefficient on RHS:

► $S := \{F \leq P \mid \mathbf{m} \text{ beyond } F\}$

$$\longrightarrow \mathbf{m} \in T_F P \iff F \notin S$$

► S is polyhedral complex with face vector (g_0, \dots, g_d)

► coefficient is

$$\sum_{i=0}^d (-1)^i (f_i - g_i) = 1 - \sum_{i=0}^d (-1)^i g_i = 1 - 1 = 0$$

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▷ recall $\Phi : \hat{L} \rightarrow L$ from summable series to rational functions

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Brion's Theorem

▷ recall $\Phi : \hat{L} \longrightarrow L$ from summable series to rational functions

▶ want to consider cones C with **lineality**: $\text{lineal } C := C \cap (-C)$

\Leftrightarrow maximal linear subspace in C

Example ▷ $P := [0, 3]$ $C^+ = [0, \infty) \subseteq \mathbb{R}$ $C^- = 3 - C^+ = (-\infty, 3]$

$$\triangleright P = C^+ + C^-$$

$$\triangleright \hat{G}_{C^+}(t) = \sum_{k \geq 0} t^k \quad \text{and} \quad \hat{G}_{C^-}(t) = \sum_{k \leq 3} t^k = t^3 \sum_{k \geq 0} t^{-k}$$

$$\triangleright G_{C^+}(t) = \frac{1}{1-t} \quad \text{and} \quad G_{C^-}(t) = t^3 \frac{1}{1-\frac{1}{t}} = \frac{t^4}{1-t}$$

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▷ this can only be true if $\Phi(\hat{G}_{\mathbb{R}}(t)) = 0!$

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► **Prop:** $C \subseteq \mathbb{R}^d$ polyhedral cone

If lineal $C \neq \{0\}$, then $\Phi(\hat{G}_C(\mathbf{t})) = 0$

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$P \subseteq \mathbb{R}^d$ rational d -polytope

Then $G_P(\mathbf{t}) = \sum_{\mathbf{v} \text{ vertex of } P} G_{T_{\mathbf{v}}P}(\mathbf{t})$

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proof: ► Apply Φ to both sides of the Brianchon-Gram identity

► only vertex cones are pointed

Brion's Theorem

Example

$$\triangleright P = \text{conv} \begin{pmatrix} 0 & 1 & 1 & 0 \\ 0 & 0 & 1 & 1 \end{pmatrix}$$



$$\begin{aligned} G_P(x, y) &= \frac{1}{(1-x)(1-y)} + \frac{x}{(1-\frac{1}{x})(1-y)} + \frac{y}{(1-x)(1-\frac{1}{y})} + \frac{xy}{(1-\frac{1}{x})(1-\frac{1}{y})} \\ &= \frac{1 + (-x^2) + (-y^2) + x^2y^2}{(1-x)(1-y)} \\ &= \frac{(1-x^2)(1-y^2)}{(1-x)(1-y)} \\ &= 1 + x + y + xy \end{aligned}$$

Geometry of Numbers

Minkowski's Theorem

▷ Λ lattice of rank d

▶ Theorem (Blichfeldt)

$S \subseteq \mathbb{R}^d$ (Lebesgue measurable) set with $\text{vol } S > \det \Lambda$

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▷ Pick $\mathbf{a} \in S_{\mathbf{x}} \cap S_{\mathbf{y}}$ and let $\mathbf{p} := \mathbf{a} + \mathbf{x}, \mathbf{q} := \mathbf{a} + \mathbf{y}$.

▷ Then $\mathbf{p}, \mathbf{q} \in S$ and $\mathbf{p} - \mathbf{q}$ is a non-zero lattice point

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$K \subseteq \mathbb{R}^d$ centrally symmetric and convex, $\text{vol } K > 2^d \det \Lambda$

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$$\frac{1}{d!} 2^d \det \Lambda \leq \lambda_1 \cdot \lambda_2 \cdots \lambda_d \cdot \text{vol } K \leq 2^d \det \Lambda$$

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▶ **Prop:** In **fixed dimension** we can find \mathbf{w} in time polynomial time in $\log^3 \max(\|\mathbf{b}_i\|)$

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$$\begin{aligned} \rightarrow \text{index } C_j &= |\det(\mathbf{v}_1, \dots, \mathbf{v}_{j-1}, \mathbf{w}, \mathbf{v}_{j+1}, \dots, \mathbf{v}_d)| \\ &= \sum_{k=1}^d |\alpha_k| \cdot |\det(\mathbf{v}_1, \dots, \mathbf{v}_{j-1}, \mathbf{v}_k, \mathbf{v}_{j+1}, \dots, \mathbf{v}_d)| \\ &= |\alpha_j| \cdot |\det(\mathbf{v}_1, \dots, \mathbf{v}_d)| \\ &= \sqrt[d]{\text{index}(C)} \text{index}(C) = \text{index}(C)^{\frac{d-1}{d}} \end{aligned}$$

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$$\varepsilon_j := \begin{cases} 0 & \text{if } \dim C_j < d \\ 1 & \text{if } \det(\mathbf{v}_1, \dots, \mathbf{v}_d) = \det(\mathbf{v}_1, \dots, \mathbf{v}_{j-1}, \mathbf{w}, \mathbf{v}_{j+1}, \dots, \mathbf{v}_d) \\ -1 & \text{otherwise.} \end{cases}$$

$$\longrightarrow \hat{G}_C(\mathbf{t}) = \sum_{j=1}^d \varepsilon_j \hat{G}_{C_j}(\mathbf{t}) + \text{lower dimensional contributions}$$

- this decomposition has
- at most d cones in dimension d
 - at most $2^d d$ cones in arbitrary dimension

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► repeat this decomposition until we reach **unimodular cones**

→ n steps such that $\left(\text{index}(C)^{\frac{d-1}{d}}\right)^n < 2$

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- ▶ **Theorem:** $d \in \mathbb{Z}_{>}$ fixed. P d -dimensional polyhedron P given in exterior description.

There is a **polynomial time** algorithm that computes $G_P(\mathbf{t})$ in the form

$$G_P(\mathbf{t}) = \sum_{i \in I} \varepsilon_i \frac{\mathbf{t}^{\mathbf{a}_i}}{(1 - \mathbf{t}^{\mathbf{v}_{i1}}) \cdots (1 - \mathbf{t}^{\mathbf{v}_{id}})}$$

where ▶ $\varepsilon_i \in \{-1, 1\}$

▶ $\mathbf{a} \in \mathbb{Z}^d$

▶ $\mathbf{v}_{ij} \in \mathbb{Z}^d - \{\mathbf{0}\}$ for all i, j ,

Short Basis

► **Theorem**(Lenstra, Lenstra, Lovász)

Λ lattice, basis $\mathbf{b}_1, \dots, \mathbf{b}_d$, $\ell := \max(\|\mathbf{b}_i\|)$.

There is $M = M(d)$ s.th. we obtain basis $\mathbf{v}_1, \dots, \mathbf{v}_d$ in time $\mathcal{O}(d^6 \log^3 \ell)$ with

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$$\text{Then } \mathbf{w} = \sum_{i=1}^d \lambda_i \mathbf{v}_i \quad \text{with} \quad |\lambda_i| \leq \sqrt{d}M$$

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- ▶ This reduces finding \mathbf{w} to a finite enumeration problem!

Theorem of Scott

Scott's Theorem

► Theorem

P lattice polygon with $i \geq 1$ interior lattice points. Then either

1. $P = 3\Delta_2$, so $\text{vol } P = \frac{9}{2}$ and $i = 1$, or
2. $\text{vol } P \leq 2(i + 1)$

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► proof

can assume ► P is contained in rectangle $\text{conv} \begin{pmatrix} 0 & p' & p' & 0 \\ 0 & 0 & p & p \end{pmatrix}$

► p minimal with this property

► $2 \leq p \leq p'$

q_b, q_t intersection with bottom and top

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→ $b \leq q_b + q_t + 2p$ and $a \geq \frac{p(q_b + q_t)}{2}$ and $a \geq \frac{b}{2}$

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► $b \leq q_b + q_t + 2p$

$$a \geq \frac{p(q_b + q_t)}{2}$$

$$a \geq \frac{b}{2}$$

► case distinction

1. $p = q_b + q_t = 3$

2. $p = 2$ or $q_b + q_t \geq 4$

3. $p = 3$ and $q_b + q_t \leq 2$

4. $p \geq 4$ and $q_b + q_t \leq 2$

h^* of polygons

► Theorem

$h_2^* t^2 + h_1^* t + 1$ for $h_1^*, h_2^* \in \mathbb{N}$ is h^* -polynomial of a lattice polygon if and only if

1. $h_2^* = 0$ and h_1^* is arbitrary. Then P has no interior lattice points.
2. $h_2^* = 1$ and $h_1^* = 7$. Then $P \cong 3\Delta_2$.
3. $1 \leq h_2^* \leq h_1^* \leq 3h_2^* + 3$. Then P has interior lattice points.

► proof

$$\triangleright h_2^* = i \quad \text{and} \quad h_1^* = b + i - 3 \quad \text{and} \quad a = \frac{1}{2}(1 + h_1^* + h_2^*)$$

$$\text{(Scott)} \quad \longrightarrow \quad \text{if } i \geq 1 \text{ and } P \neq 3\Delta_2, \text{ then } h_1^* \leq 3h_2^* + 3$$

$$b \geq 3 \quad \longrightarrow \quad \text{if } i \geq 1, \text{ then } h_2^* = i \leq b + i - 3 = h_1^*$$

this proves necessity

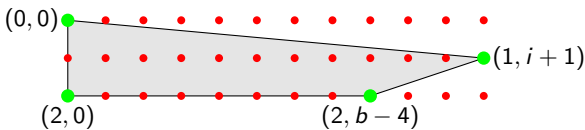
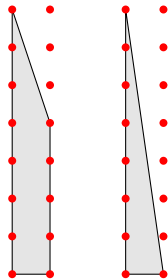
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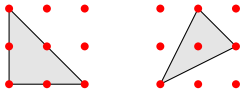
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$$h_2^* \leq h_1^* \leq 3h_2^* + 3 \iff 3 \leq b \leq 2i + 6$$



h^* of polygons

► Corollary

$c_2 t^2 + c_1 t + 1$ with $c_1, c_2 \in \frac{1}{2}\mathbb{Z}$ and $c_1 \geq \frac{3}{2}$

is Ehrhart polynomial of a lattice polygon P

if and only if one of the following three conditions is satisfied:

1. $c_1 - c_2 = 1$. Then P has no interior lattice points.
2. $c_1 = c_2 = 9/2$. Then P is $3\Delta_2$.
3. $c_1 \leq \frac{c_2}{2} + 2$. Then P has interior lattice points.